

The Role of Microeconomic Theory in Networked Systems' Performance

D. Teneketzis

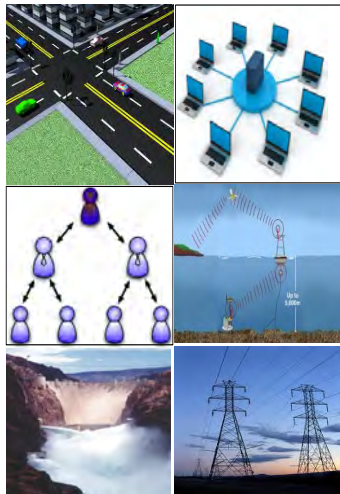
in collaboration with

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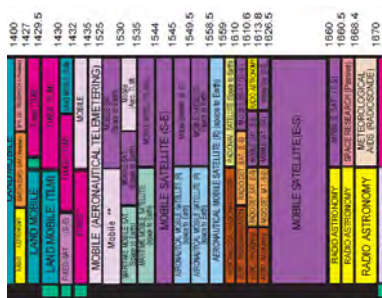
September 2013

Decentralized Systems

- Energy Markets
- Communication Networks
- Social Networks
- Auctions
- Transportation Systems
- Networked Control Systems
- Environmental Monitoring Systems
- ⋮



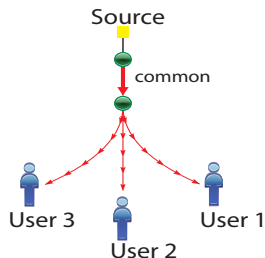
Radio Spectrum Allocation in Wireless Networks



- Radio waves are what allow cell phones to communicate.
- Radio waves travel across different frequencies of radio spectrum.
- In the U.S., every business or individual who wants to broadcast using radio waves must acquire a license from the FCC.
- If several agents use the same frequency band, they experience interference that potentially reduces their benefits.

-what is an efficient way the FCC can use to allocate radio spectrum to strategic agents?

Bandwidth Allocation in Wired Networks



- Agents **share** a link with a finite bandwidth capacity.
- Agents have different and **Private** values/utilities.

-what is the optimal (maximizing social welfare) way to allocate the finite rate/bandwidth to a set of agents?

Economic Systems: Auctions

 <p>32 GB Apple iPad 3Gs Price: \$649 Bid: \$30.67 SOLD Save: 92%</p>	 <p>Apple iMac 27" 1 TB Price: \$1553 Bid: \$85.78 SOLD Save: 95%</p>
 <p>Apple Macbook Pro 15.4" 250 GB Price: \$1285 Bid: \$65.84 SOLD Save: 75%</p>	 <p>Apple TV Black wi-fi Price: \$99 Bid: \$12.15 SOLD Save: 88%</p>
<p>Department Stores are Ripping you Off Save up to 90% Off</p> <p>QuiBids BID NOW</p> <p><small>* simulation of typical auction results</small></p>	

- Buyers in auctions make bidding decisions based on their private valuation of the object being sold.
 - *what is the optimal auction that allocates the object to the buyer whose value is the highest (maximizing revenue)?*

Energy Systems: Energy Procurement from a Strategic Energy Generator

- Strategic Buyer
 - Utility Maximizer
 - Strategic Seller
 - Conventional and renewable energy resources
 - Private generation technology
- How should the buyer negotiate and sign a contract for energy procurement with the seller?*

What are the **common** features of these problems?

Key features in decentralized resource allocation problems

- Many **agents** have to **share** the system's **limited resources**.
- Each **agent** may have **different information** about the system, and agents' decisions/actions affect the other agents' information.
- **Agents** may behave **strategically**. They exchange information with one another and determine communication and decision strategies that lead to the objective of the problem (network objective).

Mechanism Design Theory:

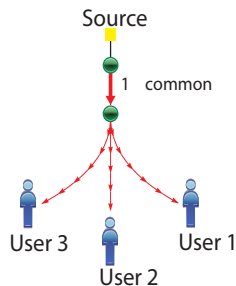
Formal treatment of
decentralized resource allocation
problems with **strategic** agents.

What is a **resource allocation problem**?
(centralized problem)

A resource allocation problem

Bandwidth allocation in wired networks:

- **Environment space:** U_1, U_2, U_3 , and the fixed network
- **Allocation Space:** Bandwidth (x_1, x_2, x_3)
- **Objective:** Maximizing social welfare



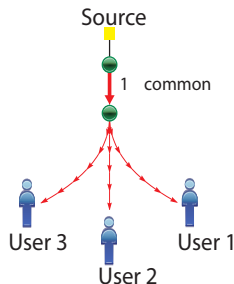
The centralized problem:

$$\begin{aligned} \max_{(x_1, x_2, x_3)} \quad & U_1(x_1) + U_2(x_2) + U_3(x_3) \\ \text{s.t.} \quad & x_1 + x_2 + x_3 \leq 1 \\ & x_1 \geq 0 \\ & x_2 \geq 0 \\ & x_3 \geq 0 \end{aligned}$$

A resource allocation problem

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The centralized problem:

$$\max_{(x_1, x_2, x_3)} U_1(x_1) + U_2(x_2) + U_3(x_3)$$

$$\text{s.t. } x_1 + x_2 + x_3 \leq 1$$

$$x_1 \geq 0$$

$$x_2 \geq 0$$

$$x_3 \geq 0$$

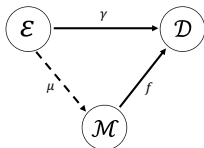
Utilities are **private** information
Agents are **strategic**

What is **Implementation Theory**?

Mechanism design

Implementation theory deals with Strategic agents.

- **Mechanism/game form:** (\mathcal{M}, f)



- **Message/Strategy space \mathcal{M} :** Set of messages agents can communicate to other agents.
 - **Outcome function f :** Determines resource allocations at each message profile.
- **Game $(\mathcal{M}, f, \{u_i\})$** is induced by game form (\mathcal{M}, f) :
Players – network agents; Strategy set – \mathcal{M} ; Players' utilities – $\{u_i(f(\mathbf{m}))\}$.
- **Equilibrium/solution concept:** Nash equilibrium, Bayesian Nash equilibrium, etc. *Appropriate solution concept determined by the information structure of the game.*

Nash equilibrium (NE):

- A message profile $\mathbf{m}_{\mathcal{N}}^* := (m_1^*, m_2^*, \dots, m_N^*)$ with property

$$u_i(f_i(\mathbf{m}_{\mathcal{N}}^*)) \geq u_i(f_i(m_i, \mathbf{m}_{\mathcal{N}/i}^*)) \quad \forall m_i \in \mathcal{M}_i, \quad \forall i \in \mathcal{N}.$$

*

Desirable properties of a **Mechanism**

Desirable properties of a mechanism/game form:

(I) Implementation in Nash equilibria: A game form (\mathcal{M}, f) “implements the goal correspondence γ in Nash equilibria” if, for all problem environments,

Set of allocations at **All Nash equilibria** \subset Set of **optimal** centralized allocations

\equiv **ALL** the **Nash equilibria** of the game induced by the game form are **efficient**.

Desirable properties of a mechanism/game form:

(II) Individual rationality: A game form (\mathcal{M}, f) is individually rational if, for all agents,

Utility at all Nash equilibria \geq Utility before/without participating in the allocation process specified by the game form

\equiv All the **agents VOLUNTARILY participate** in the game induced by the game form.

Desirable properties of a mechanism/game form:

(III) Budget balance: A game form (\mathcal{M}, f) is budget balanced if,

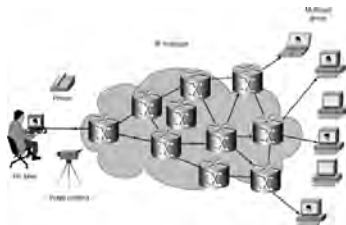
Net money transfer in the system = 0

≡ **Sum** of the **taxes** is *always* equal to **zero**.

Multi-rate Multicast service provisioning
problem with **strategic** agents

Multi-rate Multicast Technology

History



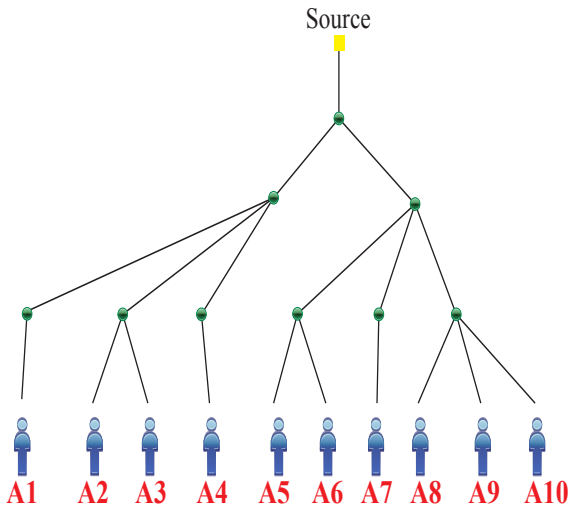
- Multi-rate Multicast technology is used as a protocol in [wired networks](#).
- Initiated in 1990, (called **MBONE**), by USC, MIT and the Lawrence Berkeley National Lab.
- **Large-scale widely used communication infrastructure**

Literature Survey

- See A. Kakhbod, D. Teneketzis, “An Efficient Game Form for Multi-rate Multicast Service Provisioning” *IEEE Journal on Selected Areas in Communication, Special Issue on the Economics of Communication Networks and Systems*, Vol. 30, No. 10, December 2012, pp. 2093-2104
- All previous literature assumes non-strategic users.

Multi-rate Multicast Technology

The main feature



Multi-rate Multicast Technology

The main feature

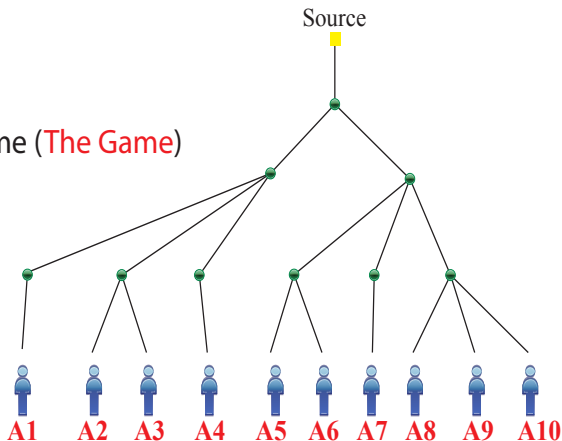
4 types of Programs:

Movie 1

Movie 2

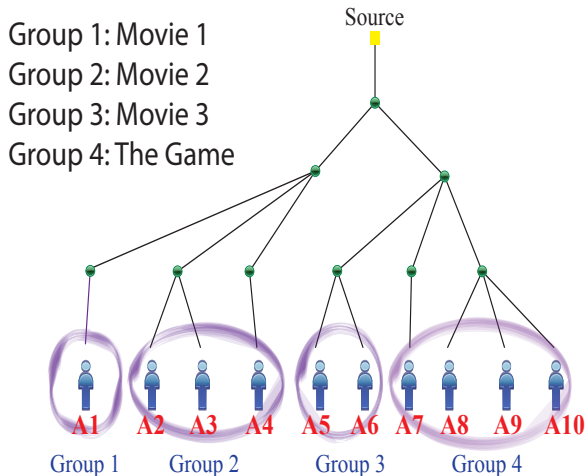
Movie 3

The super bowl game (**The Game**)



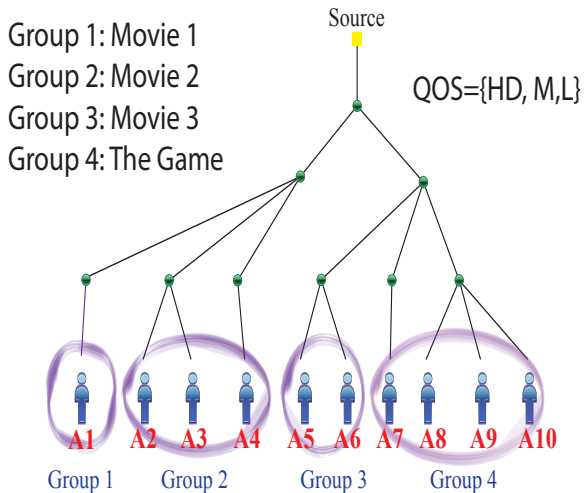
Multi-rate Multicast Technology

The main feature



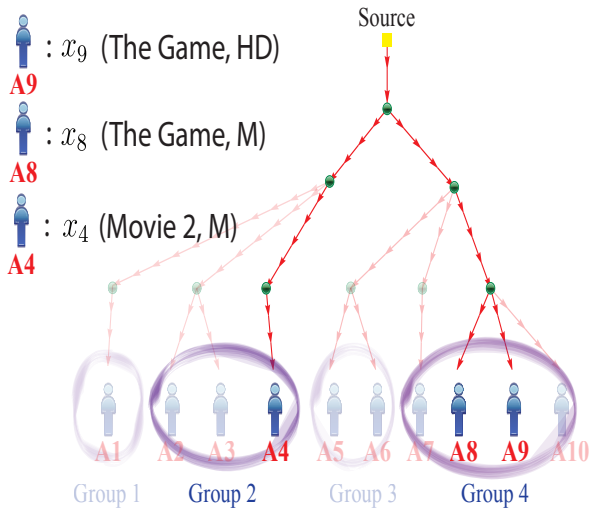
Multi-rate Multicast Technology

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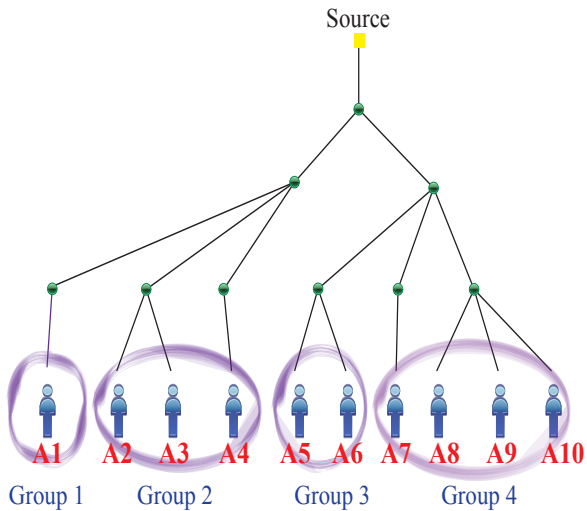
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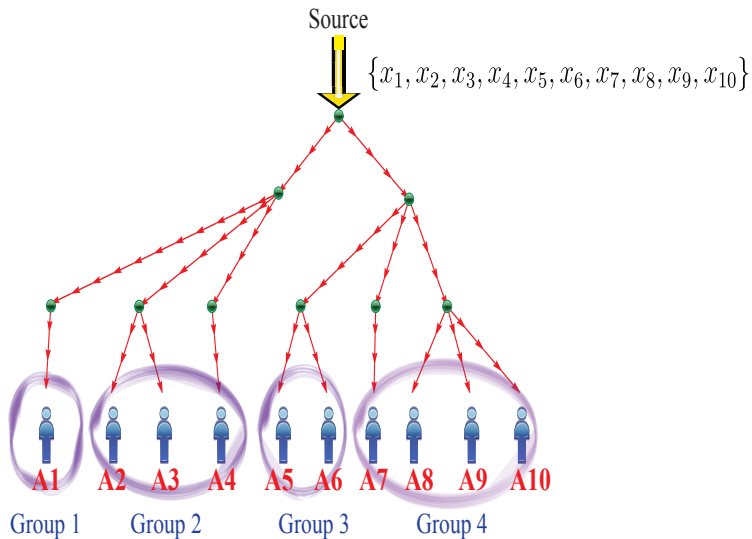
Multi-rate Multicast Technology

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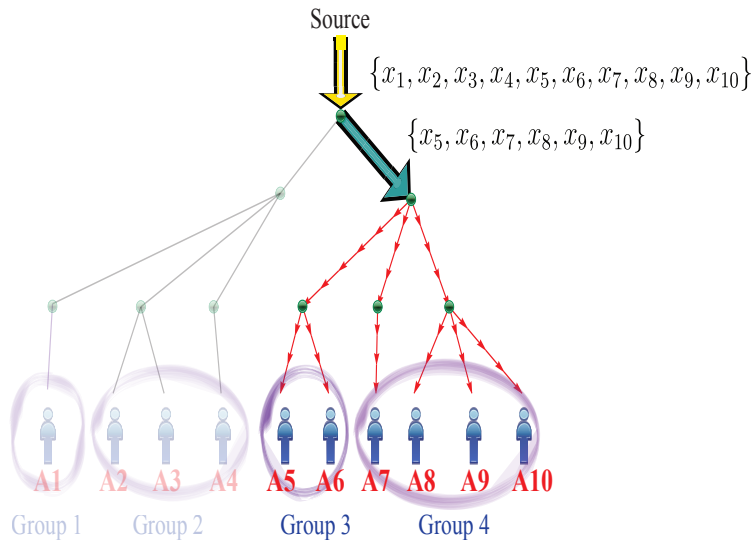
Multi-rate Multicast Technology

The main feature



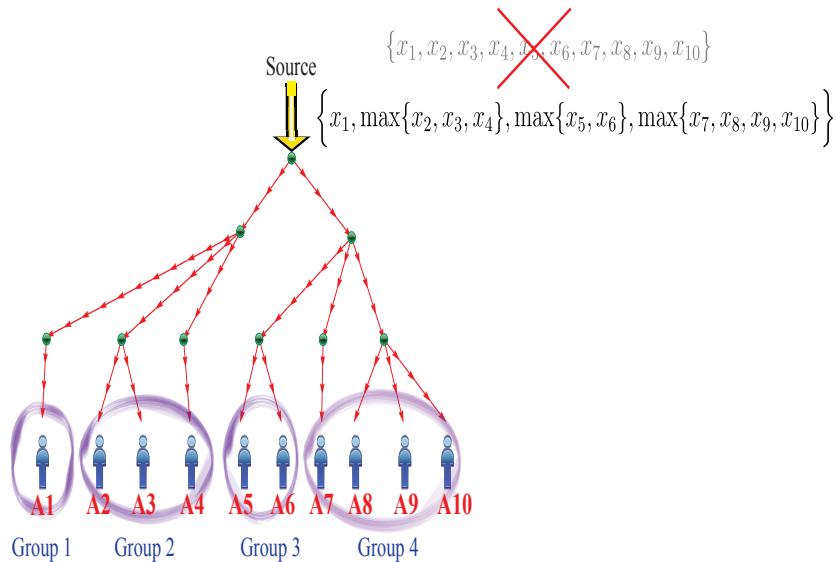
Multi-rate Multicast Technology

The main feature



Multi-rate Multicast Technology

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Multi-rate Multicast Technology

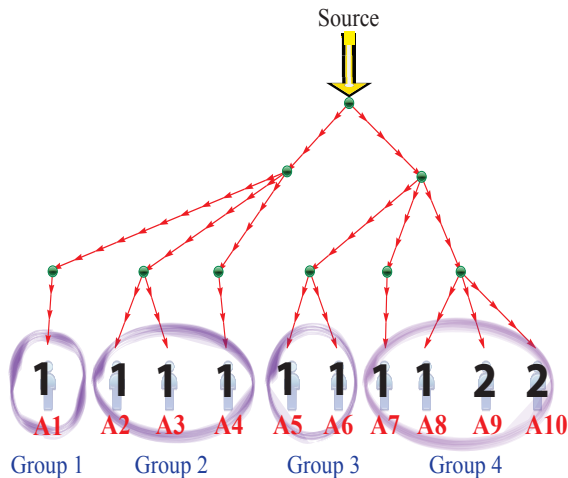
The main feature

(Movie i, M), (The Game, M) \rightarrow 1

(The Game, HD) \rightarrow 2

A_1 : (Movie1, M), A_2, A_3, A_4 : (Movie 2, M), A_5, A_6 : (Movie 3, M)

A_7, A_8 : (The Game, M), A_9, A_{10} : (The Game, HD),



Multi-rate Multicast Technology

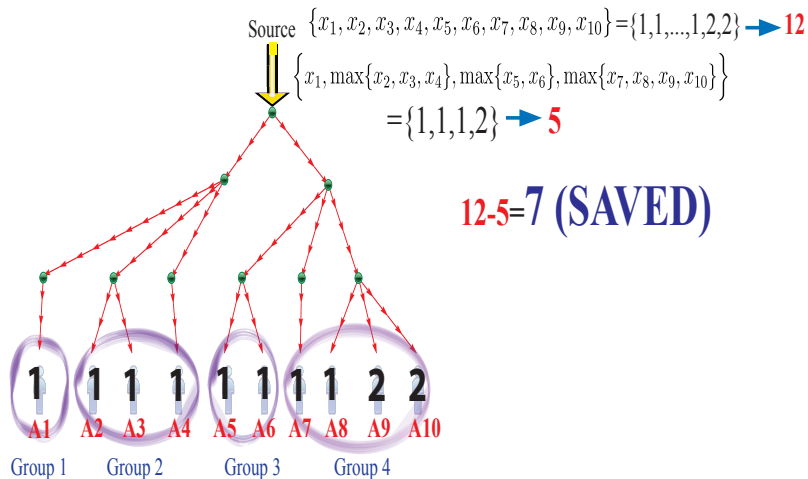
The main feature

(Movie i, M), (The Game, M) \rightarrow 1

(The Game, HD) \rightarrow 2

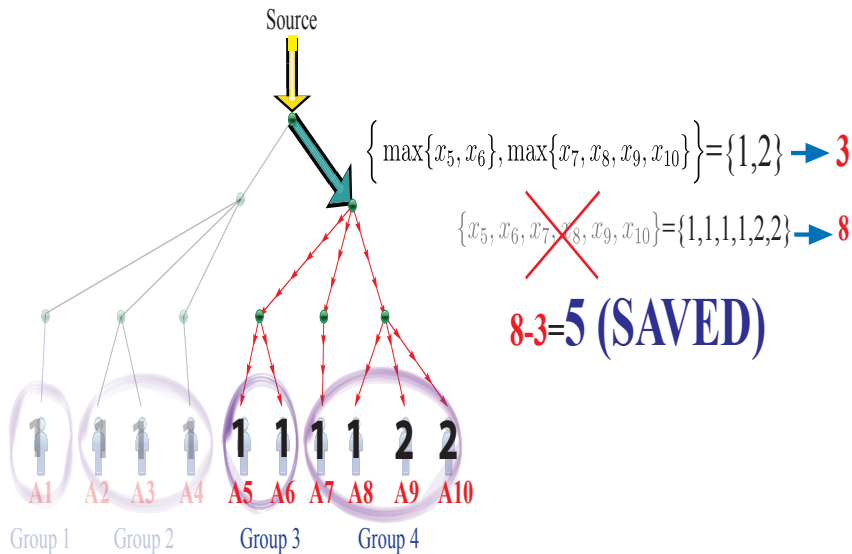
A_1 : (Movie1, M), A_2, A_3, A_4 : (Movie 2, M), A_5, A_6 : (Movie 3, M)

A_7, A_8 : (The Game, M), A_9, A_{10} : (The Game, HD),



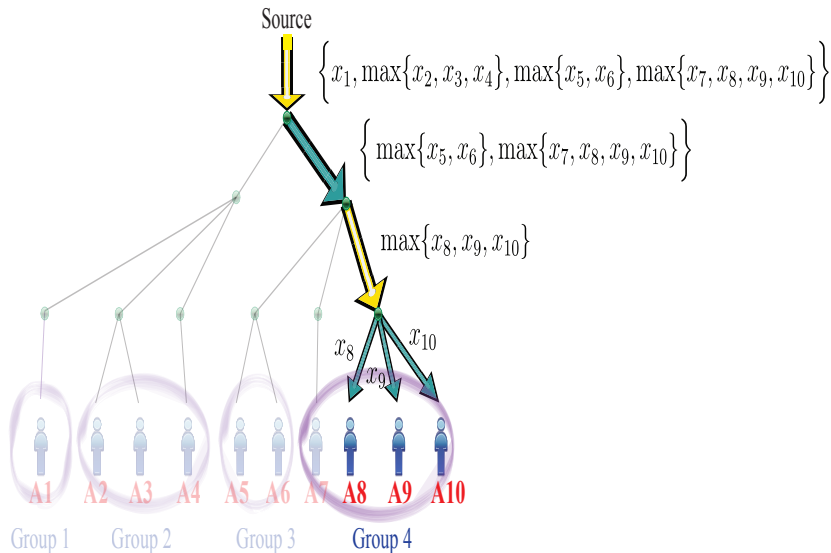
Multi-rate Multicast Technology

The main feature



Multi-rate Multicast Technology

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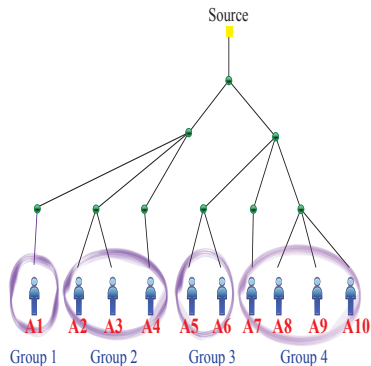


What is the **Centralized** problem?

Multi-rate Multicast Service Provisioning Problem

The Centralized Problem (P_C)

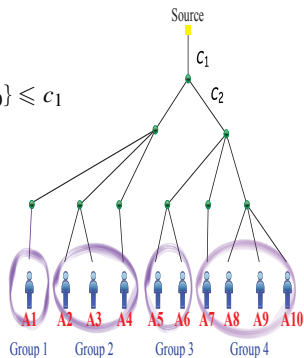
$$\begin{aligned} \max_{(x_1, x_2, \dots, x_{10})} \quad & U_1(x_1) + U_2(x_2) + \dots + U_{10}(x_{10}) \\ \text{s.t.} \quad & \text{Capacity Constraint at each link is satisfied} \\ & x_1 \geq 0 \\ & x_2 \geq 0 \\ & \vdots \\ & x_{10} \geq 0 \end{aligned}$$



Multi-rate Multicast Service Provisioning Problem

The Centralized Problem (P_C)

$$\begin{aligned} \max_{(x_1, x_2, \dots, x_{10})} \quad & U_1(x_1) + U_2(x_2) + \dots + U_{10}(x_{10}) \\ \text{s.t.} \quad & x_1 + \max\{x_2, x_3, x_4\} + \max\{x_5, x_6\} + \max\{x_7, x_8, x_9, x_{10}\} \leq c_1 \\ & \max\{x_5, x_6\} + \max\{x_7, x_8, x_9, x_{10}\} \leq c_2 \\ & \vdots \end{aligned}$$



Multi-rate Multicast Service Provisioning Problem

The Centralized Problem (P_C)

Problem (P_C)

$$\begin{aligned} \max_{(\mathbf{x}_N)} \quad & \sum_{G_i \in \mathcal{N}} \sum_{(j, G_i) \in G_i} U_{(j, G_i)}(x_{(j, G_i)}) \\ \text{s.t.} \quad & \sum_{G_i \in Q_l} \max_{(j, G_i) \in G_i(l)} x_{(j, G_i)} \leq c_l \quad \forall l \in \mathbf{L} \\ & x_{(j, G_i)} \geq 0 \quad \forall (j, G_i) \end{aligned}$$

Key **difficulties** in obtaining a
centralized optimal solution!

(1) Information is **Decentralized**.

- Every agent's utility function is its **private** information.
- Utility function of agent i , $V_i(x_i, t_i) = u_i(x_i) - t_i$
 $t_i \rightarrow$ tax

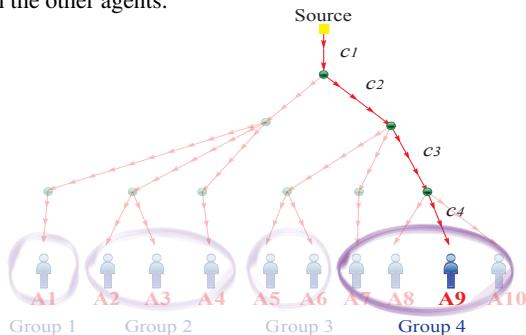
(2) Agents are **Strategic**.

Above difficulties imply that we must solve a *decentralized optimization problem with strategic agents*.

Decentralized optimization problem with strategic agents:

Problem formulation

- (I1) Each agent knows *only* his own utility. ($V_i(x_i, t_i) = u_i(x_i) - t_i$)
- (I2) Each agent behaves *strategically*. The agent's objective is to maximize his own utility function.
- (I3) The network operator knows the topology and resources of the network.
- (I4) The network operator receives requests for service and announces to each agent
 - The group to which the agent belongs.
 - The set of links that form agent's route.
 - The capacity of each link in his route.
- (I5) Each strategic agent competes for resources (bandwidth) at each link of his route with the other agents.



Problem formulation

Objective

Develop a **game form** that

- (1) is individually rational.
- (2) results in budget balance.
- (3) implements in Nash equilibria the optimal solution of Problem (P_C)
ALL the NE are EFFICIENT.

Problem (P_C)

$$\begin{aligned} \max_{(\mathbf{x}_N)} \quad & \sum_{G_i \in \mathcal{N}} \sum_{(j, G_i) \in G_i} U_{(j, G_i)}(x_{(j, G_i)}) \\ \text{s.t.} \quad & \sum_{G_i \in Q_l} \max_{(j, G_i) \in G_i(l)} x_{(j, G_i)} \leq c_l \quad \forall l \in \mathbf{L} \\ & x_{(j, G_i)} \geq 0 \quad \forall (j, G_i) \end{aligned}$$

Contribution:

Developed a **game form** that

- (1) is individually rational.
- (2) results in budget balance at equilibrium.
- (3) **implements in Nash equilibria the optimal solution of Problem (P_C)**
ALL the NE are EFFICIENT.

Problem (P_C)

$$\begin{aligned} \max_{(\mathbf{x}_N)} \quad & \sum_{G_i \in \mathcal{N}} \sum_{(j, G_i) \in G_i} U_{(j, G_i)}(x_{(j, G_i)}) \\ \text{s.t.} \quad & \sum_{G_i \in Q_l} \max_{(j, G_i) \in G_i(l)} x_{(j, G_i)} \leq c_l \quad \forall l \in \mathbf{L} \\ & x_{(j, G_i)} \geq 0 \quad \forall (j, G_i) \end{aligned}$$

What are the main features of the
game form/mechanism?

Multi-rate Multicast Service Provisioning with Strategic agents

Key features of the Game form

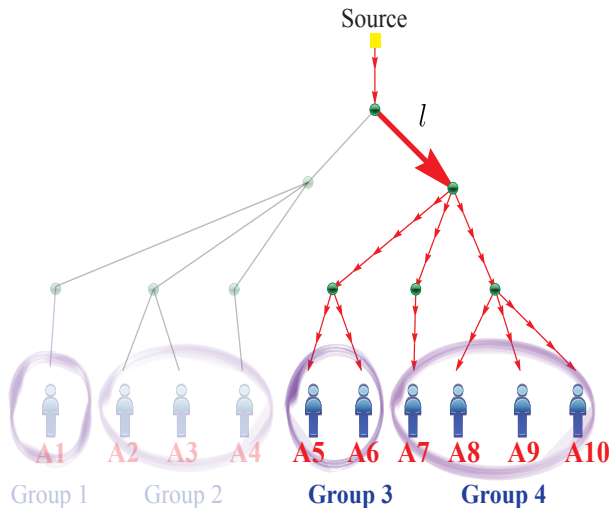
Multi-rate Multicast service provisioning with strategic agents is the combination of a Market problem and a Public good problem

- The resource allocation among groups is a market problem.
- The resource allocation among the agents of the same group is a public good problem.

Multi-rate Multicast Service Provisioning with Strategic agents

Key features of the Game form

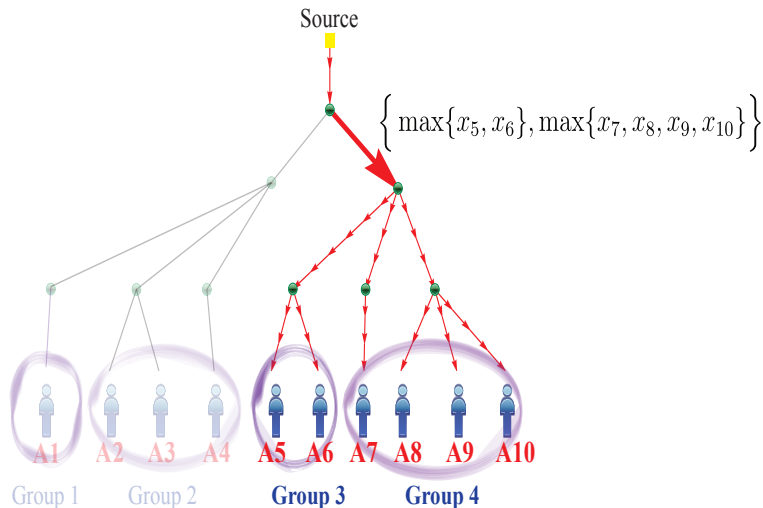
The Market component:



Multi-rate Multicast Service Provisioning with Strategic agents

Key features of the Game form

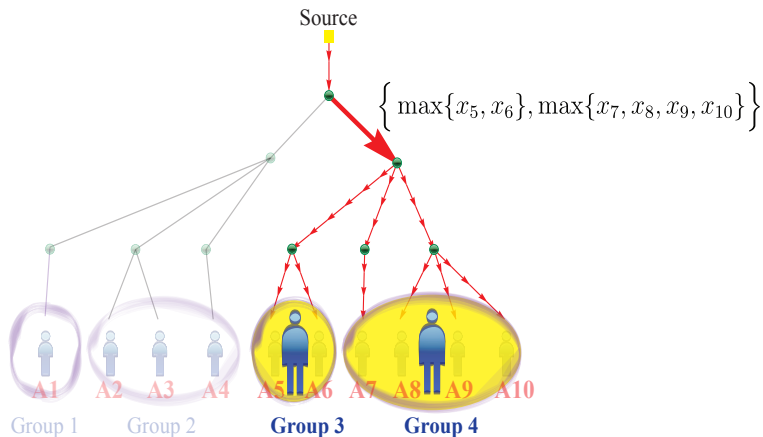
The Market component:



Multi-rate Multicast Service Provisioning with Strategic agents

Key features of the Game form

The Market component:



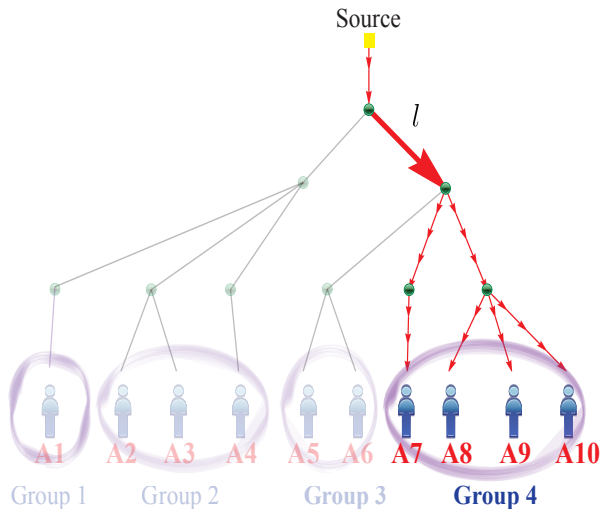
→ **A. Kakhbod** and D. Teneketzis, “An efficient game form for unicast service provisioning problem,”

IEEE Transaction on Automatic Control (TAC), vol 57, no. 2, pp. 392 - 404, 2012.

Multi-rate Multicast Service Provisioning with Strategic agents

Key features of the Game form

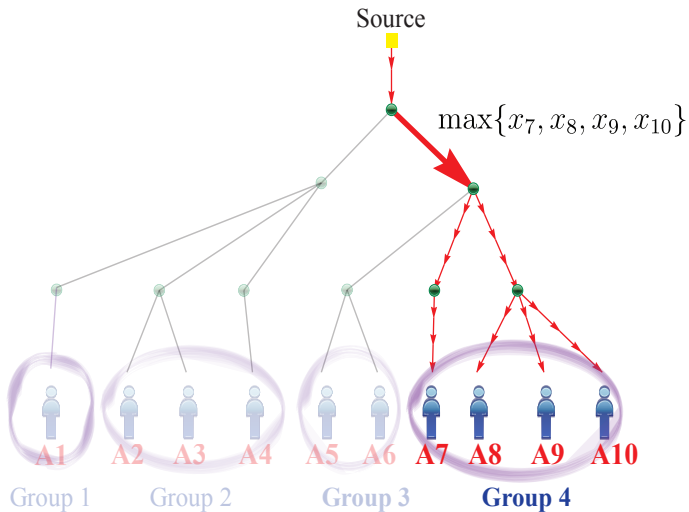
The Public good component:



Multi-rate Multicast Service Provisioning with Strategic agents

Key features of the Game form

The Public good component:



Specification of
 $M :=$ Message Space and $f :=$ Outcome function

Specification of the game form

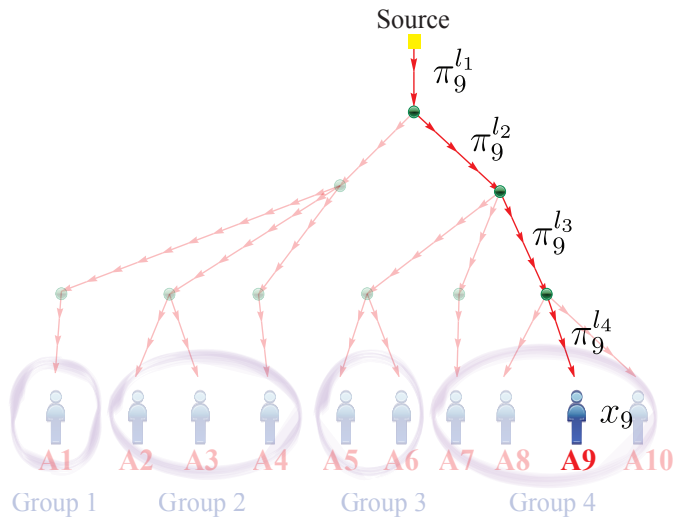
Message Space: $M := \bigotimes_{G_i \in \mathcal{N}} \bigotimes_{(j, G_i) \in G_i} M_{(j, G_i)}$

$$\mathbf{m}_{(j, G_i)} = \left(x_{(j, G_i)}, \pi_{(j, G_i)}^{l_{j_1}}, \pi_{(j, G_i)}^{l_{j_2}}, \dots, \pi_{(j, G_i)}^{l_{j_{|R_{(j, G_i)}|}}} \right)$$

- $x_{(j, G_i)}$:= **bandwidth/rate** agent (j, G_i) requests
- $\pi_{(j, G_i)}^{l_{j_k}}$:= **price per unit of bandwidth** agent (j, G_i) is willing to pay at link l_{j_k} of his route.

Specification of the game form

Message Space: $M := \bigotimes_{G_i \in \mathcal{N}} \bigotimes_{(j, G_i) \in G_i} M_{(j, G_i)}$



Specification of the game form

Outcome function: f

For any message profile \mathbf{m} , $\mathbf{m} \in M$,

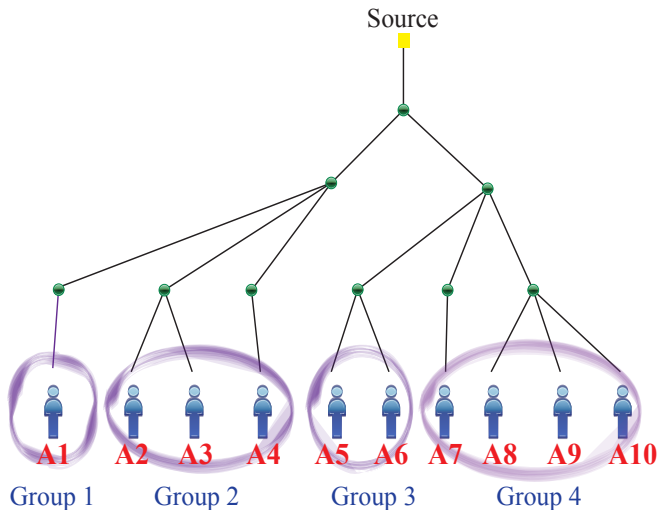
$$f(\mathbf{m}) = [(\text{bandwidth}, \text{tax})_{(j, G_i)}, \dots] \quad \text{bandwidth and tax for each agent}$$

For each agent

The tax (subsidy) is defined at **each link** of its route based on the number of groups using the link.

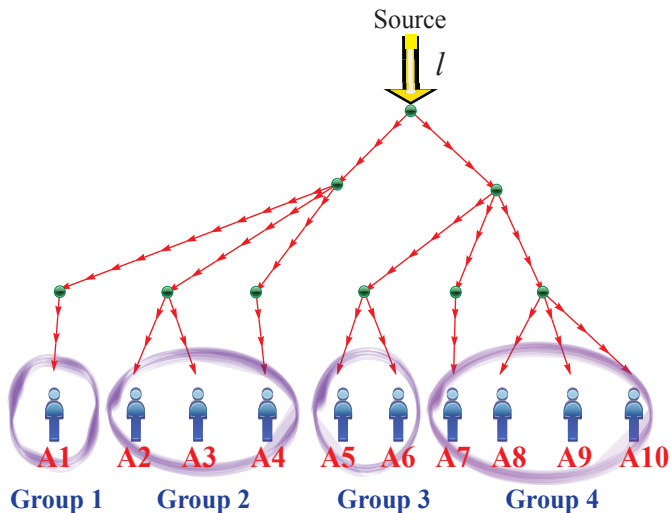
Specification of the game form

Outcome function: f



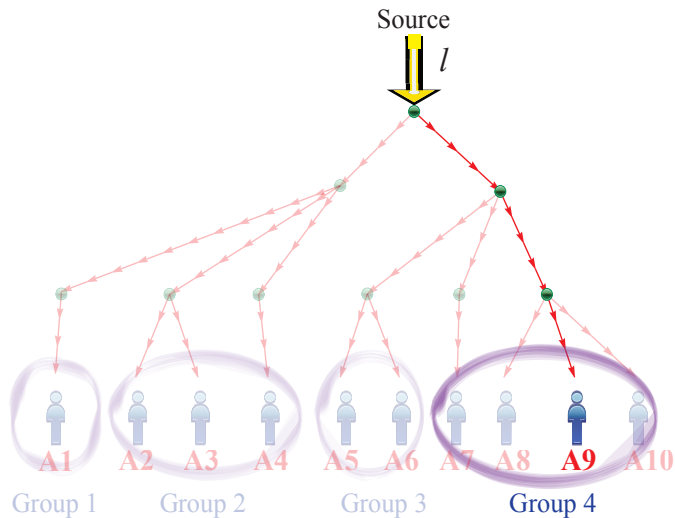
Specification of the game form

Outcome function: f



Specification of the game form

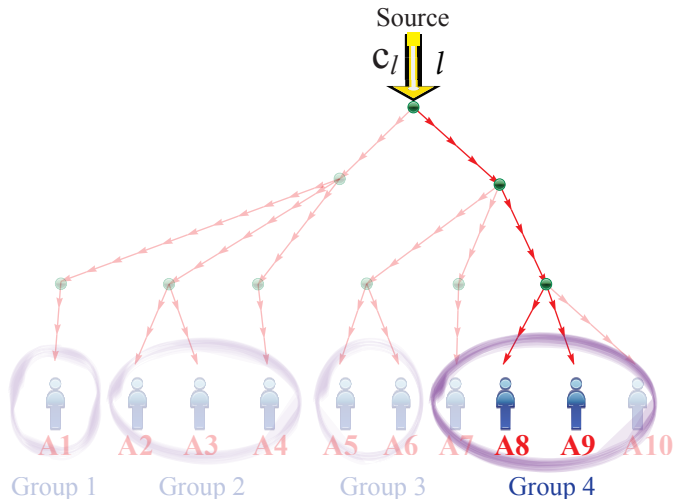
Outcome function: f



Specification of the game form

Outcome function: f

A₈ and **A₉** request the **maximum** bandwidth in **G₄** at link l



Interpretation of the tax function

Public good terms, Market terms

$$t_{A9}^l = \left[\Delta_2^{A9}(l) + \Delta_3^{A9}(l) + \Delta_4^{A9}(l) \right] \times \Delta_1^{A9}(l).$$

$$\Delta_1^{A9}(l) := \mathbb{I}\{x_9 \text{ is the the maximum request in } G_4 \text{ at link } l\}$$

$$\Delta_2^{A9}(l) := \pi_8^l x_9$$

$$\Delta_3^{A9}(l) := \frac{\left(P_{G_4(l)} - P_{-G_4(l)} - \eta_{+}^l \right)^2}{2} - P_{-G_4(l)} \left(P_{G_4(l)} - P_{-G_4(l)} \right) \left[\frac{\mathcal{E}_{-G_4(l)} + x_9}{\beta} \right]$$

$$\Delta_4^{A9}(l) := \Gamma_{G_4}^l \quad \text{(Budget Balancing term)}$$

$$\mathcal{E}_{-G_4(l)} := x_{G_1}(l) + x_{G_2}(l) + x_{G_3}(l) - c_l$$

$$\eta_{+}^l := \max\left\{0, \frac{x_{G_1}(l) + x_{G_2}(l) + x_{G_3}(l) + x_{G_4}(l) - c_l}{\gamma}\right\}$$

$$P_{-G_4}(l) := \frac{1}{3} \left(P_{G_1}(l) + P_{G_2}(l) + P_{G_3}(l) \right),$$

Specification of the game form

Outcome function: f

If A_8 requests the **maximum** bandwidth in G_4 at link l , and A_9 does **not** request the **maximum** bandwidth at l :

$$t_{A_9}^l = \pi_8 (\mathcal{E}_{-G_4(l)} + x_{G_4}(l)) (1 - \Delta_1^{A_9}(l)).$$

Properties of the **Mechanism**

[Existence and Feasibility]

- There exists at least one **pure** NE for the game induced by the game form.
- If \mathbf{m}^* is a *NE* point of the game induced by the game form, then the allocation \mathbf{x}^* is a feasible solution of Problem (\mathbf{P}_C).

[Properties of NE]

Let m^* be a NE of the game induced by game form. Then for every $l \in \mathbf{L}$ we have,

- **Each group** using the link bids the the **SAME** price = $P_{G(l)}^*$
- $P_{G(l)}^* \left[\frac{c_l - \text{Sum of the demands requested by the groups}}{\beta} \right] = 0.$

6

[Budget Balance]

The proposed mechanism/game form is always balanced budget at every allocation corresponding to NE messages.

6

[Individually Rationality]

The specified mechanism is **individually rational at all NE**, i.e., at every NE, the corresponding allocation is weakly preferred by all agents to the initial allocation.

$$\mathbf{U}_{(j, G_i)}(x_{(j, G_i)}^*) - t_{(j, G_i)}(\mathbf{m}^*) \geq 0, \quad \forall (j, G_i)$$

‘

[Nash Implementation]

Consider the allocation $(f(\mathbf{m}^*) = (\mathbf{x}^*, \mathbf{t}^*))$ corresponding to **any NE** message \mathbf{m}^* .
Then \mathbf{x}^* is an **optimal solution** of the
centralized problem (P_C).

Problem (P_C)

$$\begin{aligned} \max_{(\mathbf{x}_N)} \quad & \sum_{G_i \in \mathcal{N}} \sum_{(j, G_i) \in G_i} U_{(j, G_i)}(x_{(j, G_i)}) \\ \text{s.t.} \quad & \sum_{G_i \in Q_l} \max_{(j, G_i) \in G_i(l)} x_{(j, G_i)} \leq c_l \quad \forall l \in \mathbf{L} \\ & x_{(j, G_i)} \geq 0 \quad \forall (j, G_i) \end{aligned}$$

Conclusion

- Addressed the Multi-rate Multicast service provisioning problem for the **first** time with **strategic** agents.
- Designed a **game form** that
 - (1) is **individually rational**.
 - (2) results in **budget balance**.
 - (3) **All** the NE of the game induced by the game form are **efficient**.
- The proposed **game form** is the **first** game form that captures the features of *a Market problem* and *a Public good problem*, simultaneously.

Open Problems

- Tatonement processes for determination of NE.
- Dynamic resource allocation problems

References

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