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Local Convergence of an Incremental Algorithm for Subspace Identification

✧ Incremental Gradient

- ✧ When a cost function can be written as a sum of costs on “data blocks,” Incremental gradient performs cost function optimization one “data block” at a time.
- ✧ Great for real-time or big data applications.
- ✧ Convergence rates are poor within a local region of the solution, as compared to steepest descent or second-order methods.

✧ Manifold Optimization

- ✧ When a non-linear constraint set can be written as a Riemannian manifold, we can use manifold methods for optimization.
- ✧ Convergence results require armijo step which sometimes adds a large computational burden.

✧ Incremental Gradient

- ✧ When a cost function can be written as a sum of costs on “data blocks,” Incremental gradient performs cost function optimization one “data block” at a time.

Consider a least-squares problem of the form

$$\text{minimize}_x f(x) = \sum_{i=1}^n \|g_i(x)\|^2 .$$

✧ Incremental Gradient

$$\text{minimize}_x f(x) = \sum_{i=1}^n \|g_i(x)\|^2 .$$

Now consider the same problem but where $g_i(x)$ is a linear function of data block i , $i = 1, \dots, m$ and the incremental gradient algorithm given by [Bertsekas 99, p116] with step size α_k at iteration k . Let x^* be the optimal solution corresponding to this problem. Then:

1. There exists $\bar{\alpha} > 0$ such that if α_k is equal to some constant $\alpha \in (0, \bar{\alpha}]$ for all k , the sequence x_k converges to some vector $x(\alpha)$. Furthermore, the error $\|x_k - x(\alpha)\|$ converges to 0 linearly. Finally, we have $\lim_{\alpha \rightarrow 0} x(\alpha) = x^*$.
2. If $\alpha_k > 0$ for all k , and

$$\alpha_k \rightarrow 0, \quad \sum_{k=0}^{\infty} \alpha_k = \infty ,$$

then $\{x_k\}$ converges to x^* .

✧ Optimization on Manifolds

Consider any optimization problem on a Riemannian manifold \mathcal{M} with a retraction given from the tangent space of \mathcal{M} to \mathcal{M} .

Perform any gradient-related descent algorithm using the Armijo step size on a manifold [Absil, Mahony, Sepulchre 08, p62].

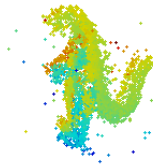
Then every limit point of the sequence of iterates is a critical point of the cost function; i.e. $\nabla f = 0$.

- ✧ Subspace Tracking with Missing Data
- ✧ GROUSE algorithm convergence rate in the full-data case
- ✧ GROUSE algorithm convergence rate with missing data
- ✧ Equivalence of grouse to a kind of missing-data incremental SVD

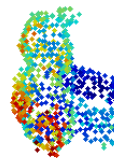
Applications that use Subspaces of R^n



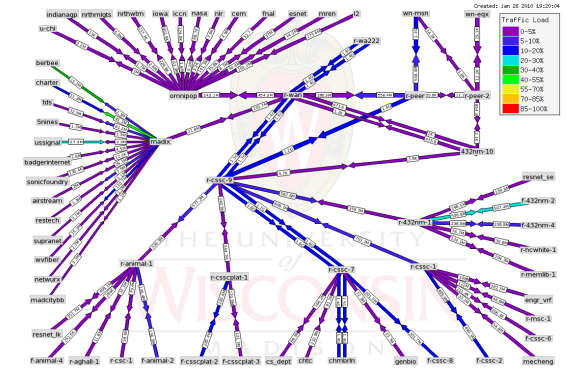
(a) Dinosaur



(b) Teddy Bear

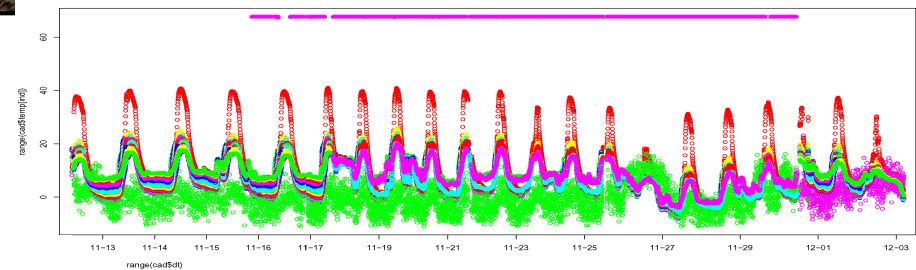


3D object modeling: when points are matched across frames, they lie in a 3D subspace.



Network data analysis: due to network connectivity constraining the flows, traffic data lie in a low dimensional subspace

Ranking based on human assessment: people's preferences have been demonstrated to lie near a low-dimensional manifold; we are using a handful of factors only

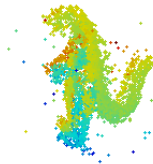


Sensor network data analysis: very spatially correlated data lie near a low-dimensional subspace

These Applications all have Missing Data



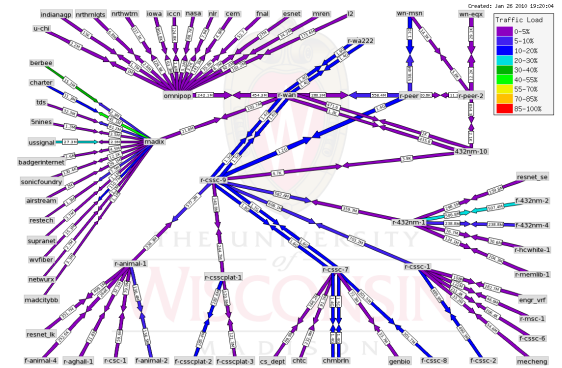
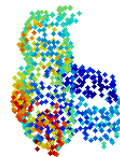
(a) Dinosaur



3D object modeling:
missing data due to obstruction from different camera angles

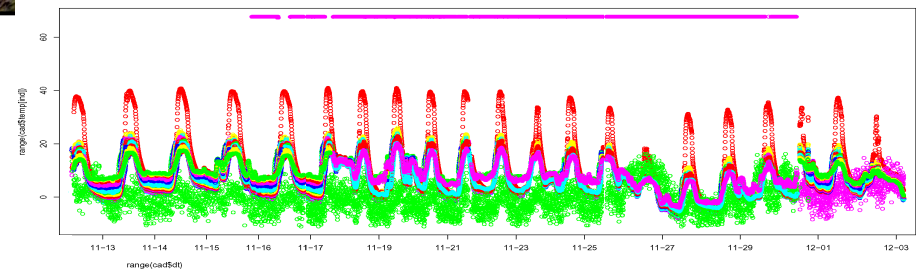


(b) Teddy Bear



Network data analysis:
missing data due to massive throughput

Ranking based on human assessment:
missing data due to impossibility of considering all alternatives

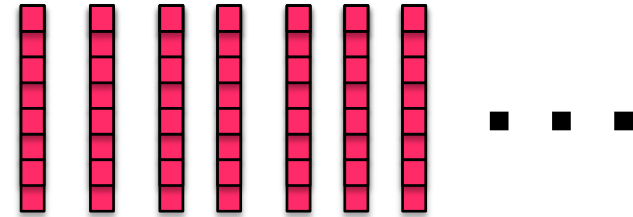


Sensor network data analysis: missing data due to cheap sensors and crummy communication links

Subspace Identification: Full Data

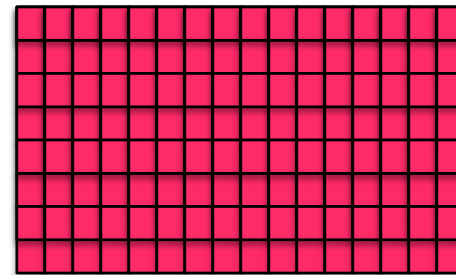
Suppose we receive a sequence of length- n vectors that lie in a d -dimensional subspace S :

$$v_1, v_2, \dots, v_t, \dots, \in S \subset \mathbb{R}^n$$



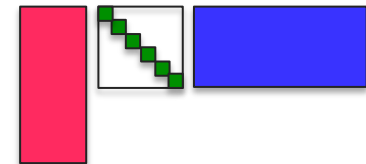
And then we collect T of these vectors into a matrix,

$$X = \begin{bmatrix} | & | & \dots & | \\ v_1 & v_2 & \dots & v_T \\ | & | & \dots & | \end{bmatrix}$$



If S is static, we can identify it as the column space of this matrix by performing the SVD:

$$X = U \Sigma V^T .$$

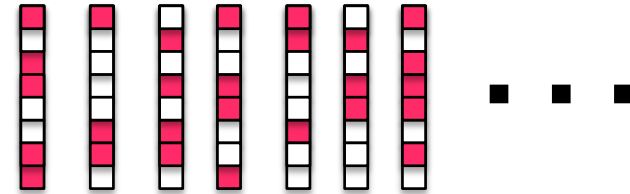


The orthogonal columns of U span the subspace S .

Subspace Identification: Missing Data

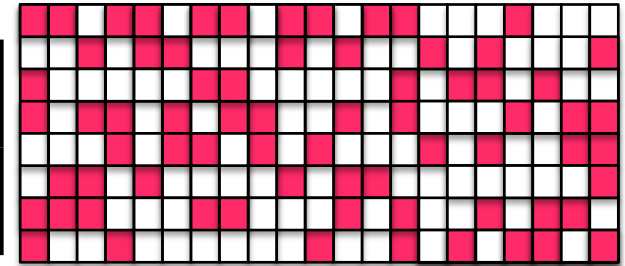
Suppose we receive a sequence of incomplete length- n vectors that lie in a d -dimensional subspace S , and $\Omega_t \subset \{1, \dots, n\}$ refers to the observed indices:

$$[v_1]_{\Omega_1}, [v_2]_{\Omega_2}, \dots, [v_t]_{\Omega_t}, \dots, \in S \subset \mathbb{R}^n$$



And then we collect T of these vectors into a matrix:

$$X = \begin{bmatrix} | & | & \dots & | \\ [v_1]_{\Omega_1} & [v_2]_{\Omega_2} & \dots & [v_T]_{\Omega_T} \\ | & | & \dots & | \end{bmatrix}$$



~~If S is static, we can identify it as the column space of this matrix by performing the SVD:~~

$$X = U\Sigma V^T .$$

The orthogonal columns of U span the subspace S .

- Seek subspace $S \subset \mathbb{R}^n$ of known dimension $d \ll n$.
- Know certain components $\Omega_t \subset \{1, 2, \dots, n\}$ of vectors $v_t \in S$, $t = 1, 2, \dots$ — the subvector $[v_t]_{\Omega_t}$.
- Assume that S is incoherent w.r.t. the coordinate directions.

We'll also assume for purposes of analysis that

- $v_t = \bar{U}s_t$, where \bar{U} is an $n \times d$ orthonormal spanning S and the components of $s_t \in \mathbb{R}^d$ are i.i.d. normal with mean 0.
- Sample set Ω_t is independent for each t with $|\Omega_t| \geq q$, for some q between d and n .
- Observation subvectors $[v_t]_{\Omega_t}$ contain no noise.

We take an incremental gradient approach to minimizing over \mathcal{S} the function

$$F(\mathcal{S}) = \sum_{i=1}^T \|[v_i - P_{\mathcal{S}}v_i]_{\Omega_i}\|_2^2.$$

Since the variable is a subspace we optimize on the Grassmannian.

Given current estimate U_t and partial data vector $[v_t]_{\Omega_t}$, where $v_t = \bar{U}s_t$:

$$w_t := \arg \min_w \|[U_t w - v_t]_{\Omega_t}\|_2^2;$$

$$p_t := U_t w_t;$$

$$[r_t]_{\Omega_t} := [v_t - U_t w_t]_{\Omega_t}; \quad [r_t]_{\Omega_t^c} := 0;$$

$$\sigma_t := \|r_t\| \|p_t\|;$$

Choose $\eta_t > 0$;

$$U_{t+1} := U_t + \left[(\cos \sigma_t \eta_t - 1) \frac{p_t}{\|p_t\|} + \sin \sigma_t \eta_t \frac{r_t}{\|r_t\|} \right] \frac{w_t^T}{\|w_t\|};$$

We focus on the (locally acceptable) choice

$$\eta_t = \frac{1}{\sigma_t} \arcsin \frac{\|r_t\|}{\|p_t\|}, \quad \text{which yields } \sigma_t \eta_t = \arcsin \frac{\|r_t\|}{\|p_t\|} \approx \frac{\|r_t\|}{\|p_t\|}.$$



To measure the discrepancy between the current estimate span(U_t) and \mathcal{S} , we use the angles between the two subspaces. There are d angles between two d -dimensional subspaces, and we call them $\phi_{t,i}$, $i = 1, \dots, d$, where

$$\cos \phi_{t,i} = \sigma_i(U_t^T \bar{U}) ,$$

where σ_i denotes the i^{th} singular value. Define

$$\epsilon_t := \sum_{i=1}^d \phi_{t,i} = d - \sum_{i=1}^d \sigma_i(U_t^T \bar{U})^2 = d - \|U_t^T \bar{U}\|_F^2 .$$

We seek a bound for $\mathbb{E}[\epsilon_{t+1} | \epsilon_t]$, where the expectation is taken over the random vector s_t for which $v_t = \bar{U} s_t$.

- ✧ Subspace Tracking with Missing Data
- ✧ **GROUSE** algorithm convergence rate in the full-data case
- ✧ GROUSE algorithm convergence rate with missing data
- ✧ Equivalence of grouse to a kind of missing-data incremental SVD

Full-data case **vastly simpler** to analyze than the general case. Define

- $\theta_t := \arccos(\|p_t\|/\|v_t\|)$ is the angle between $R(U_t)$ and \mathcal{S} that is revealed by the update vector v_t ;
- Define $A_t := U_t^T \bar{U}$, $d \times d$, nearly orthogonal when $R(U_t) \approx \mathcal{S}$. We have $\epsilon_t = d - \|A_t\|_F^2$.

Lemma

$$\epsilon_t - \epsilon_{t+1} = \frac{\sin(\sigma_t \eta_t) \sin(2\theta_t - \sigma_t \eta_t)}{\sin^2 \theta_t} \left(1 - \frac{s_t^T A_t^T A_t A_t^T A_t s_t}{s_t^T A_t^T A_t s_t} \right),$$

The right-hand side is nonnegative for $\sigma_t \eta_t \in (0, 2\theta_t)$, and zero if $v_t \in R(U_t) = \mathcal{S}_t$ or $v_t \perp \mathcal{S}_t$.

Theorem

Suppose that $\epsilon_t \leq \bar{\epsilon}$ for some $\bar{\epsilon} \in (0, 1/3)$. Then

$$E[\epsilon_{t+1} | \epsilon_t] \leq \left(1 - \left(\frac{1 - 3\bar{\epsilon}}{1 - \bar{\epsilon}}\right) \frac{1}{d}\right) \epsilon_t.$$

Since the sequence $\{\epsilon_t\}$ is decreasing, by the earlier lemma, we have $\epsilon_t \downarrow 0$ with probability 1 when started with $\epsilon_0 \leq \bar{\epsilon}$.

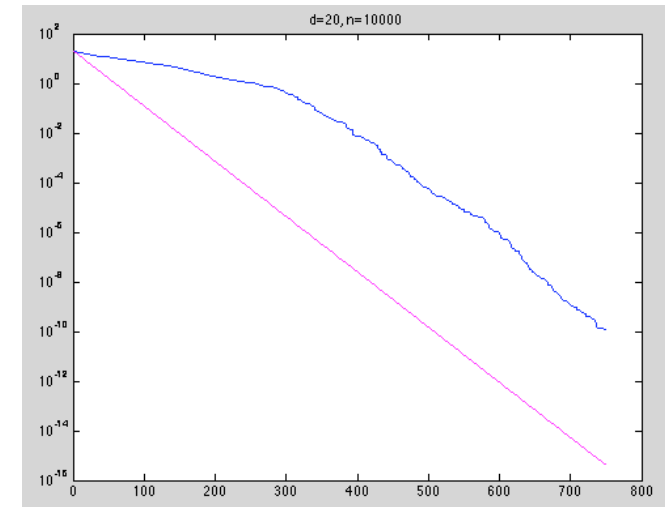
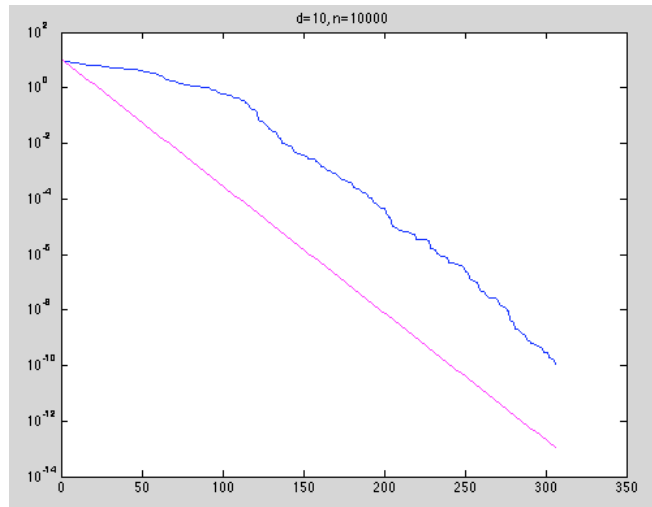
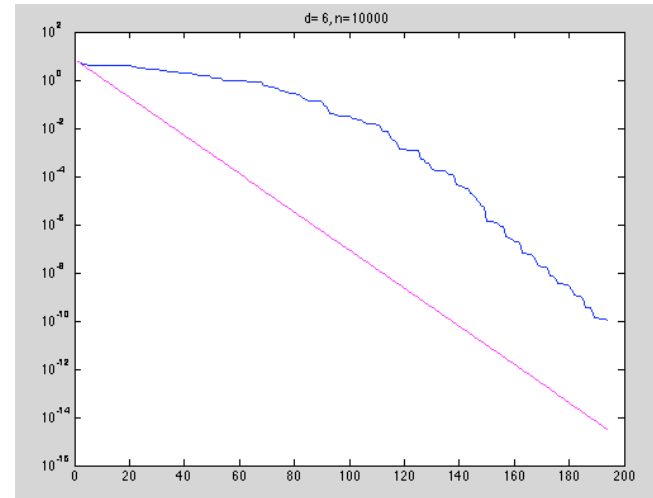
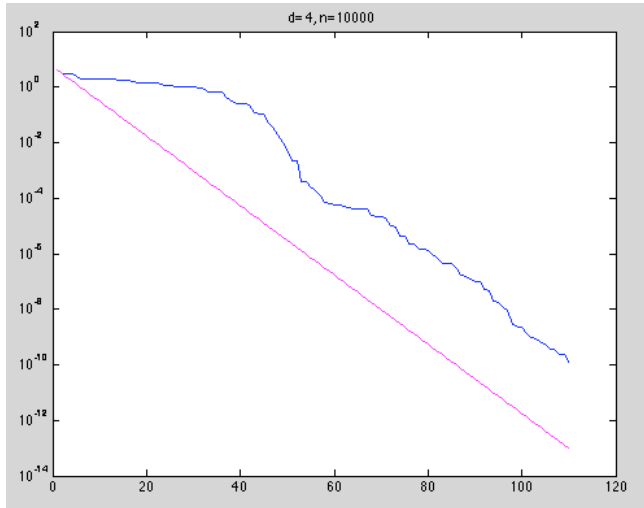
Linear convergence rate is asymptotically $1 - 1/d$.

- For $d = 1$, get near-convergence in one step (thankfully!)
- Generally, in d steps (which is number of steps to get the exact solution using SVD), improvement factor is

$$(1 - 1/d)^d < \frac{1}{e}.$$

ϵ_t versus $1-1/d$

$n=10000$
 $d=4, 6,$
 $10, 20$



- ✧ Subspace Tracking with Missing Data
- ✧ GROUSE algorithm convergence rate in the full-data case
- ✧ GROUSE algorithm convergence rate with missing data**
- ✧ Equivalence of grouse to a kind of missing-data incremental SVD

Recall, n is the ambient dimension, d the inherent dimension, we have $|\Omega| > q$ samples per vector. We have assumptions on the number of samples, the coherence in the subspaces and in the residual vectors, and we require that these assumptions hold with probability $1 - \delta$ for $\delta \in (0, .6)$. Then for

$$\epsilon_t \leq (8 \times 10^{-6})(.6 - \delta)^2 \frac{q^3}{n^3 d^2}$$

we have

$$\mathbb{E}[\epsilon_{t+1} | \epsilon_t] \leq \left(1 - (.16)(.6 - \delta) \frac{q}{nd}\right) \epsilon_t .$$

$$\epsilon_t \leq (8 \times 10^{-6})(.6 - \delta)^2 \frac{q^3}{n^3 d^2}$$

$$\mathbb{E}[\epsilon_{t+1} | \epsilon_t] \leq \left(1 - (.16)(.6 - \delta) \frac{q}{nd}\right) \epsilon_t .$$

The decrease constant is not too far from that observed in practice; we see a factor of about

$$1 - X \frac{q}{nd}$$

where X is not much less than 1.

The threshold condition on ϵ_t , however, is quite pessimistic. Linear convergence behavior is seen at much higher values.

- ✧ GROUSE algorithm convergence rate in the full-data case
- ✧ GROUSE algorithm convergence rate with missing data
- ✧ **Equivalence of grouse to a kind of missing-data incremental SVD**

Algorithm 2 iSVD: Full Data

Given U_0 , an arbitrary $n \times d$ orthonormal matrix, with $0 < d < n$; Σ_0 , a $d \times d$ diagonal matrix of zeros which will later hold the singular values, and V_0 , an arbitrary $n \times d$ orthonormal matrix.

for $t = 0, 1, 2, \dots$ **do**

 Take the current data column vector v_t ;

 Define $w_t := \arg \min_w \|U_t w - v\|_2^2 = U_t^T v_t$;

 Define

$$p_t := U_t w_t; \quad r_t := v_t - p_t;$$

 Noting that

$$\begin{bmatrix} U_t \Sigma_t V_t^T & v_t \end{bmatrix} = \begin{bmatrix} U_t & \frac{r_t}{\|r_t\|} \end{bmatrix} \begin{bmatrix} \Sigma_t & w_t \\ 0 & \|r_t\| \end{bmatrix} \begin{bmatrix} V_t & 0 \\ 0 & 1 \end{bmatrix}^T,$$

we compute the SVD of the update matrix:

$$\begin{bmatrix} \Sigma_t & w_t \\ 0 & \|r_t\| \end{bmatrix} = \hat{U} \hat{\Sigma} \hat{V}^T,$$

and set

$$U_{t+1} := \begin{bmatrix} U_t & \frac{r_t}{\|r_t\|} \end{bmatrix} \hat{U}, \quad \Sigma_{t+1} = \hat{\Sigma}, \quad V_{t+1} = \begin{bmatrix} V_t & 0 \\ 0 & 1 \end{bmatrix} \hat{V}.$$

end for

✧ We could put zeros into the matrix

- ✧ Very interesting recent results from Sourav Chatterjee on one-step “Universal Singular Value Thresholding” show that zero-filling followed by SVD reaches the minimax lower bound on MSE.
- ✧ But in the average case, we see that convergence of the zero-filled SVD is very very slow.

✧ Let's instead replace the missing entries with our prediction using the existing model

Algorithm 4 iSVD: Partial Data, Forget singular values

Given U_0 , an $n \times d$ orthonormal matrix, with $0 < d < n$;

for $t = 0, 1, 2, \dots$ **do**

 Take Ω_t and v_{Ω_t} from (2.1);

 Define $w_t := \arg \min_w \|U_{\Omega_t} w - v_{\Omega_t}\|_2^2$;

 Define vectors \tilde{v}_t, p_t, r_t :

$$(\tilde{v}_t)_i := \begin{cases} v_i & i \in \Omega_t \\ (U_t w_t)_i & i \in \Omega_t^C \end{cases}; \quad p_t := U_t w_t; \quad r_t := \tilde{v}_t - p_t;$$

Noting that

$$\begin{bmatrix} U_t & \tilde{v}_t \end{bmatrix} = \begin{bmatrix} U_t & \frac{r_t}{\|r_t\|} \end{bmatrix} \begin{bmatrix} I & w_t \\ 0 & \|r_t\| \end{bmatrix},$$

we compute the SVD of the update matrix:

$$\begin{bmatrix} I & w_t \\ 0 & \|r_t\| \end{bmatrix} = \tilde{U} \tilde{\Sigma} \tilde{V}^T,$$

and set $U_{t+1} := \begin{bmatrix} U_t & \frac{r_t}{\|r_t\|} \end{bmatrix} \tilde{U}_{:,1:d} W_t$, where W_t is an arbitrary $d \times d$ orthogonal matrix.

end for

Theorem

Suppose we have the same U_t and $[v_t]_{\Omega_t}$ at the t -th iterations of iSVD and GROUSE. Then there exists $\eta_t > 0$ in GROUSE such that the next iterates U_{t+1} of both algorithms are identical, to within an orthogonal transformation by the $d \times d$ matrix

$$W_t := \left[\frac{w_t}{\|w_t\|} \mid Z_t \right],$$

where Z_t is a $d \times (d - 1)$ orthonormal matrix whose columns span $N(w_t^T)$.

The precise values for which GROUSE and iSVD are identical are:

$$\lambda = \frac{1}{2} \left[(\|w_t\|^2 + \|r_t\|^2 + 1) + \sqrt{(\|w_t\|^2 + \|r_t\|^2 + 1)^2 - 4\|r_t\|^2} \right]$$

$$\beta = \frac{\|r_t\|^2 \|w_t\|^2}{\|r_t\|^2 \|w_t\|^2 + (\lambda - \|r_t\|^2)^2}$$

$$\eta_t = \frac{1}{\sigma_t} \arcsin \beta.$$

- ✧ Apply GROUSE analysis to ell-1 version, GRASTA
- ✧ Re-think the proof from new angles.
 - ✧ We see convergence at higher ϵ .
 - ✧ We see monotonic decrease at any random initialization.
 - ✧ We see convergence even without incoherence (but good steps are only made when the samples align).

Thank you!

Questions?